KASPER OSWALD
Ingenieure für innovative Sicherheitslösungen

RUB

Open, Sesame!

On the Security of Electronic Locks

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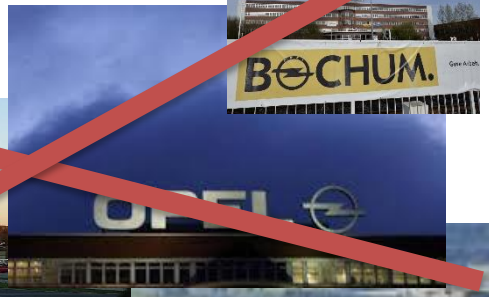
Ruhr-Uni Bochum / Kasper & Oswald

No, I did not do all this stuff alone

- Christof Paar
- Timo Kasper
- Benedikt Driessen
- Simon Küppers
- Gregor Leander
- Amir Moradi
- Ingo von Maurich
- Falk Schellenberg
- Daehyun Strobel

Ruhr-University Bochum: beautiful.



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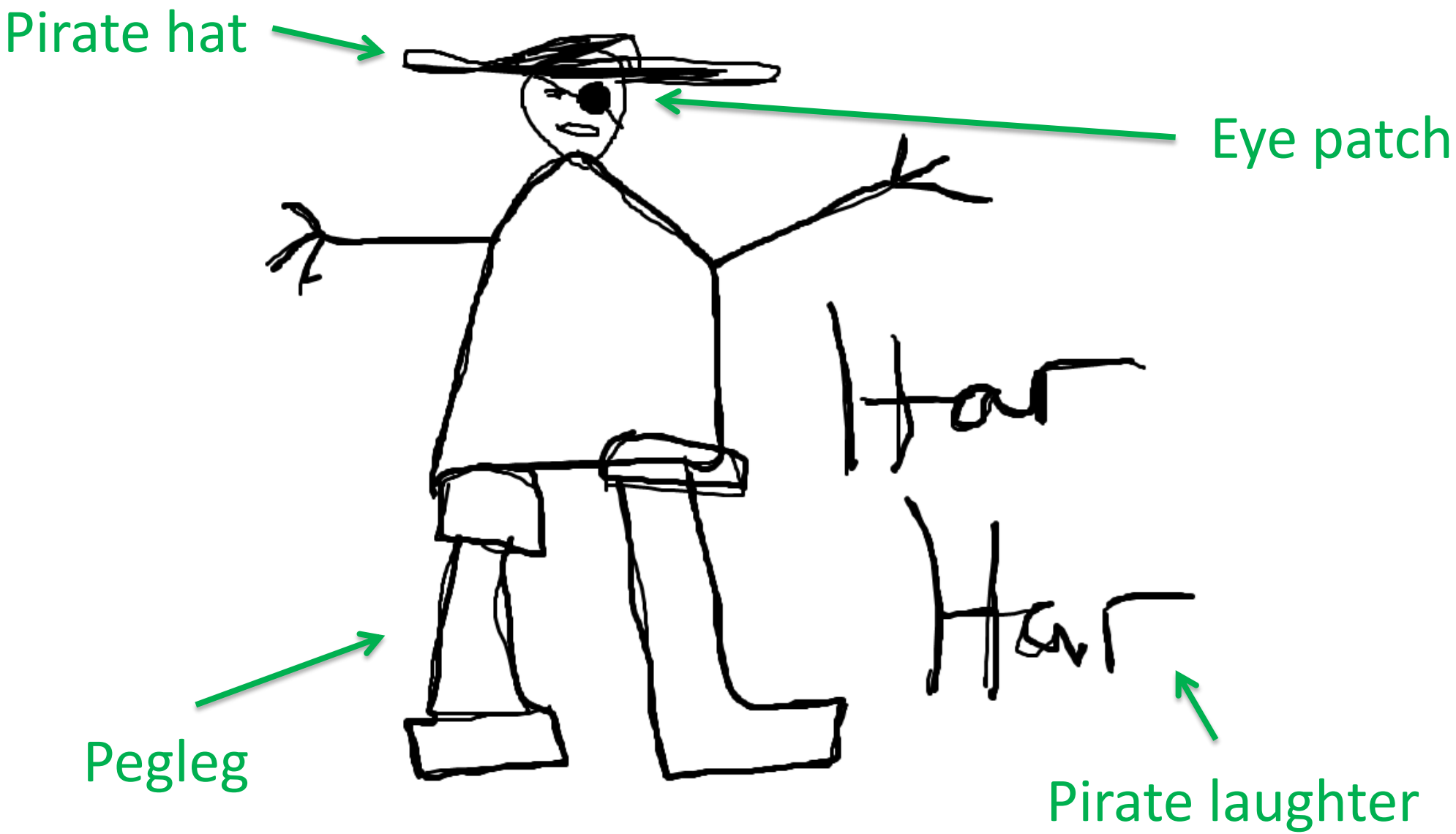
KAOS

(The life of) a typical pirate

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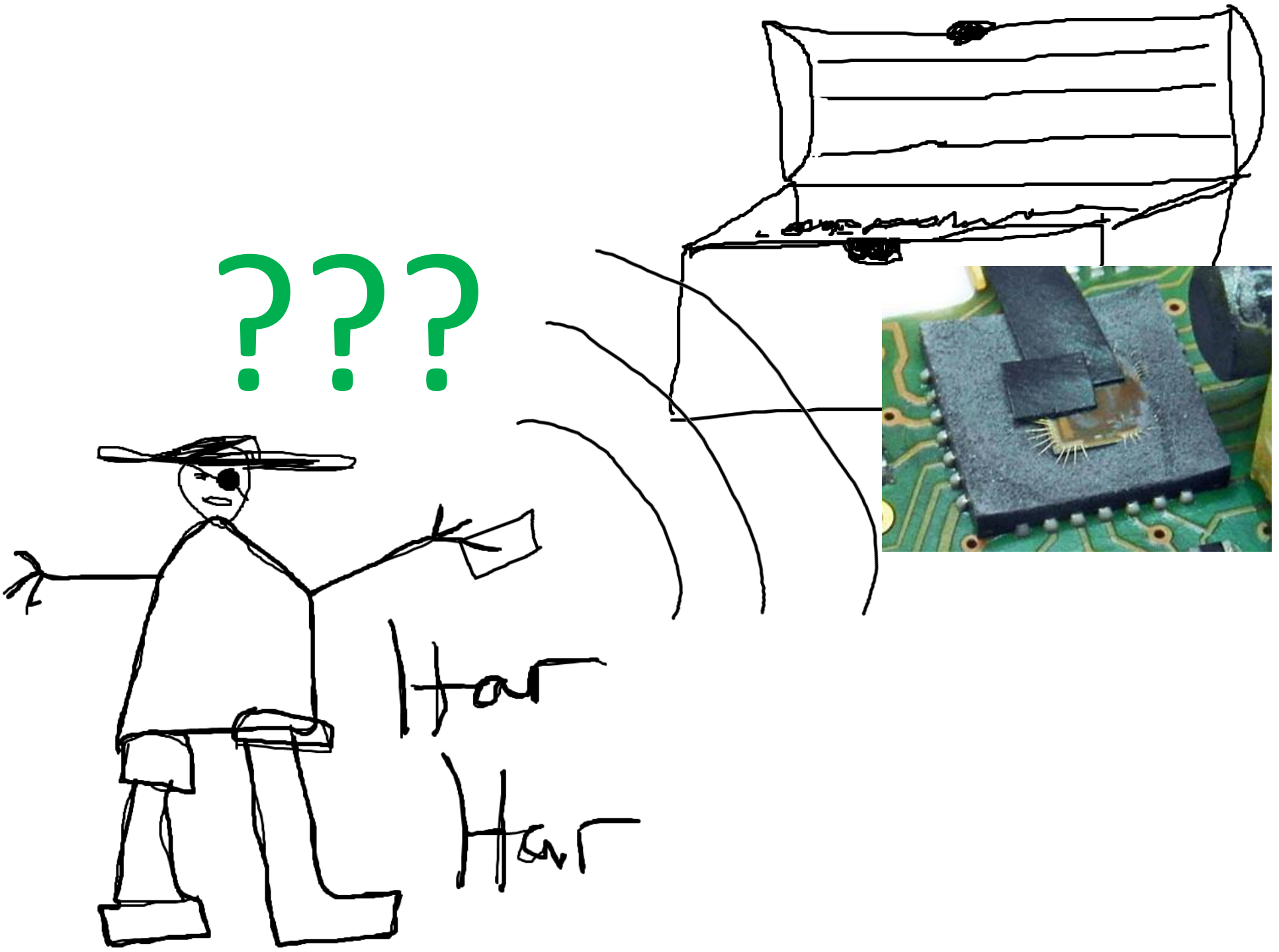
Har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Har
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„Opening“ doors – LEVEL 1

Opening doors – LEVEL 2

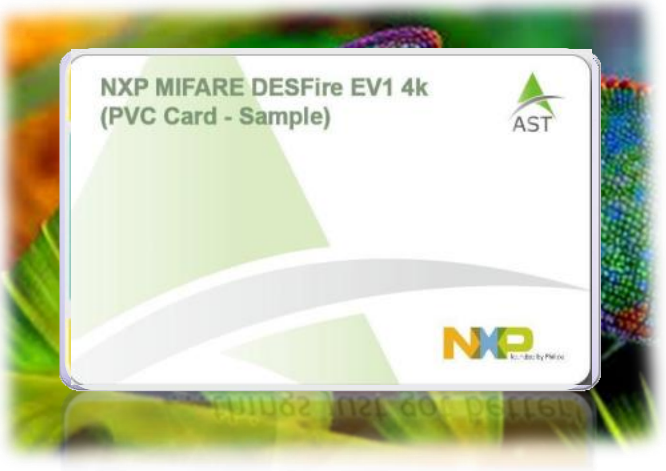
Access Control System

- Mifare Classic cards unlock doors and elevators
- Secret keys are default (0xA0A1A2A3A4A5)
- Identification by UID and
  1st block of 1st sector
- **UID usually not changeable …**

Clone on Blank Card Fails (wrong UID)

ChameleonMini

- Chameleon emulates everything *including* UID
- **Open source project:**
  https://github.com/emsec/ChameleonMini
- **Buy / Kickstarter info:**
  http://kasper-oswald.de/gb/chameleonmini



Chameleon Succeeds

(emulates everything including UID)

Quite old prototype, was actually stolen …

Level 2: Summary

- Many locks still use UID only
  (from 125 kHz to DESFire EV1…)
- Mifare Ultralight (no crypto) e.g. used for hotel rooms
- Mifare Classic (broken in 2009) still wide-spread
- Backwards compatibility & mixed systems …

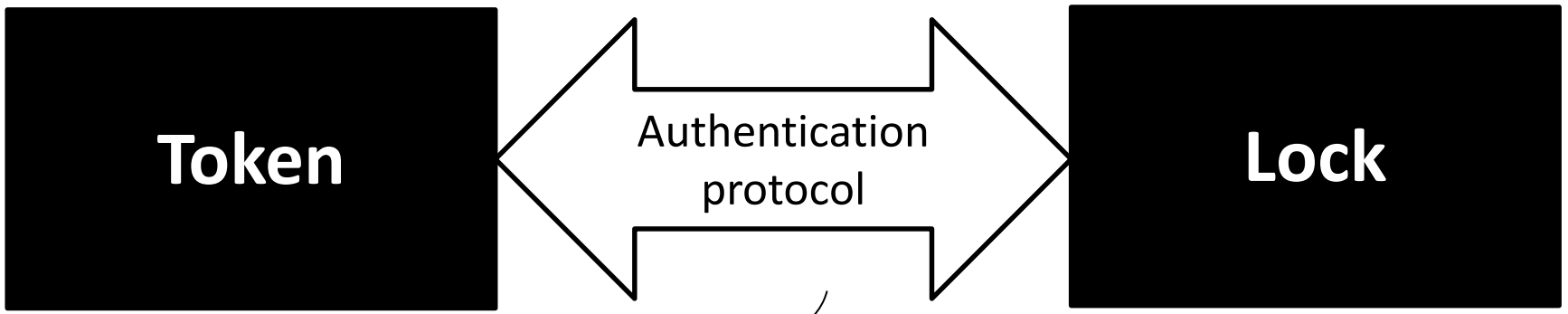
Opening doors – LEVEL 3

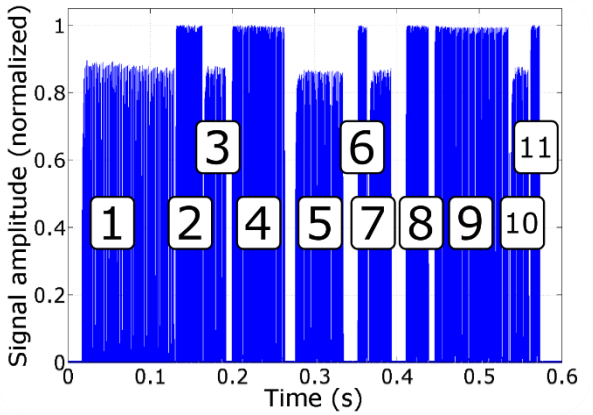
Electronic Locking System

Reverse-Engineering (1)

Black-box analysis:

Token and lock perform authentication protocol





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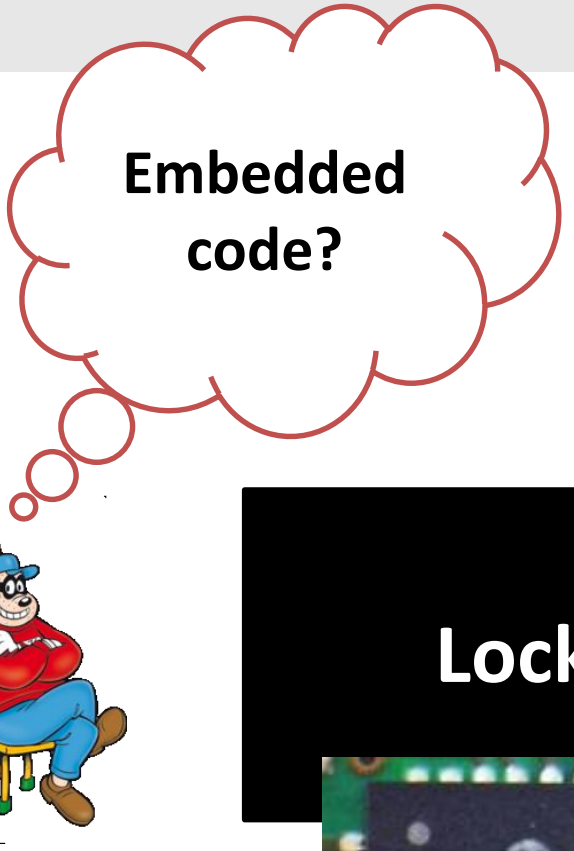
Reverse-Engineering (2)

Token

Lock

Reverse-Engineering (3)

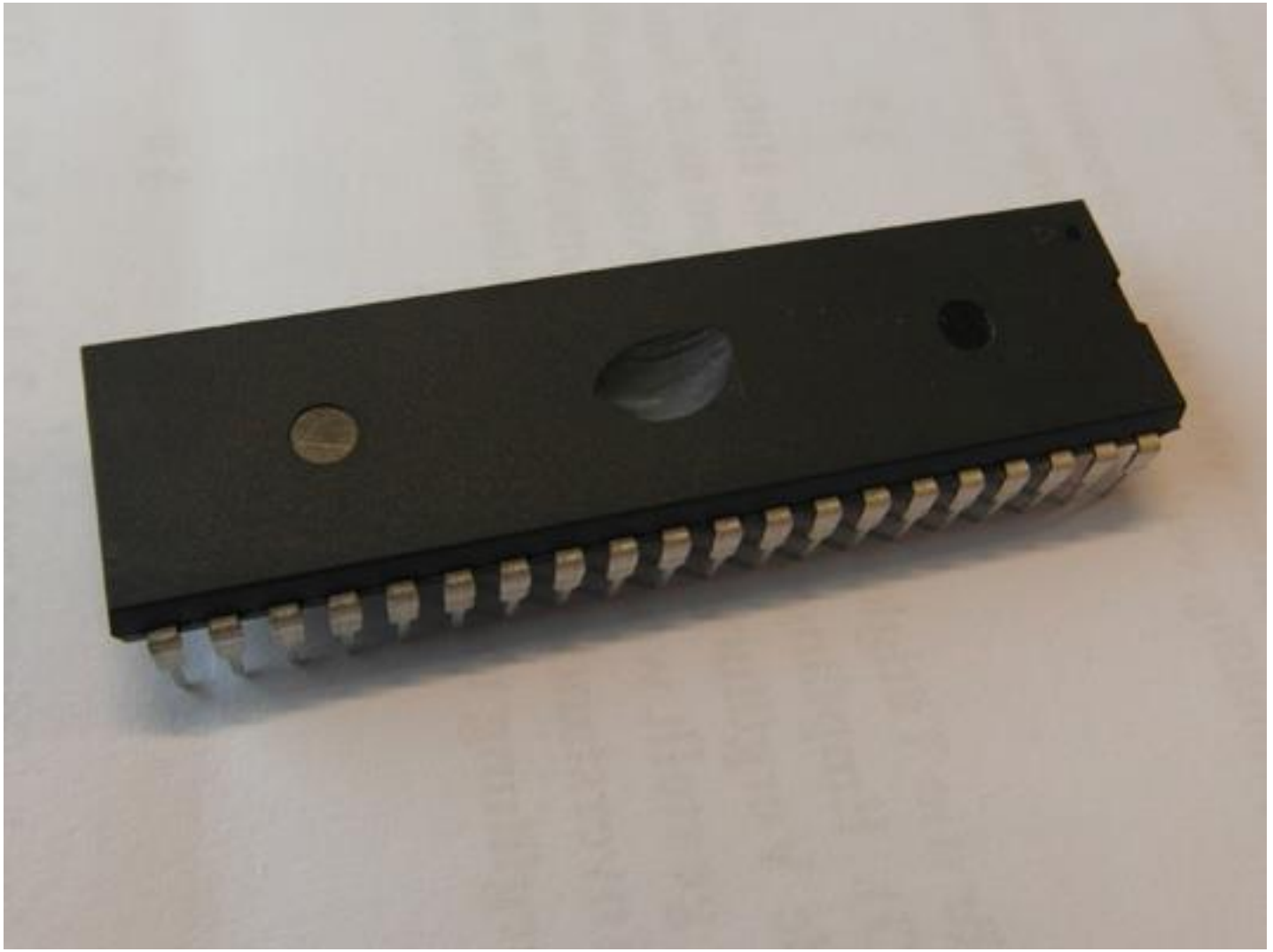


Token

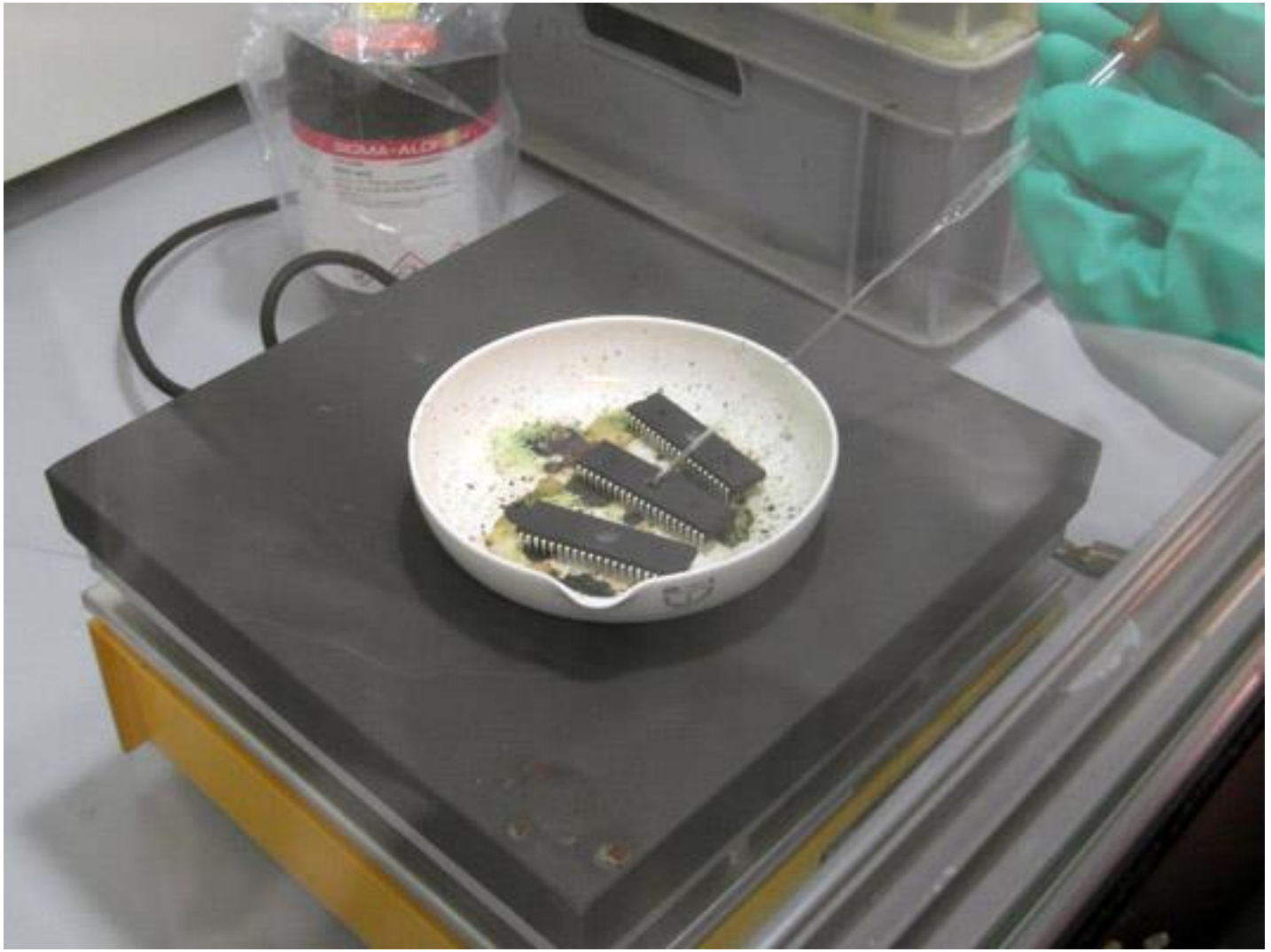
Lock

Read-out protection!

Decapping an IC (1)

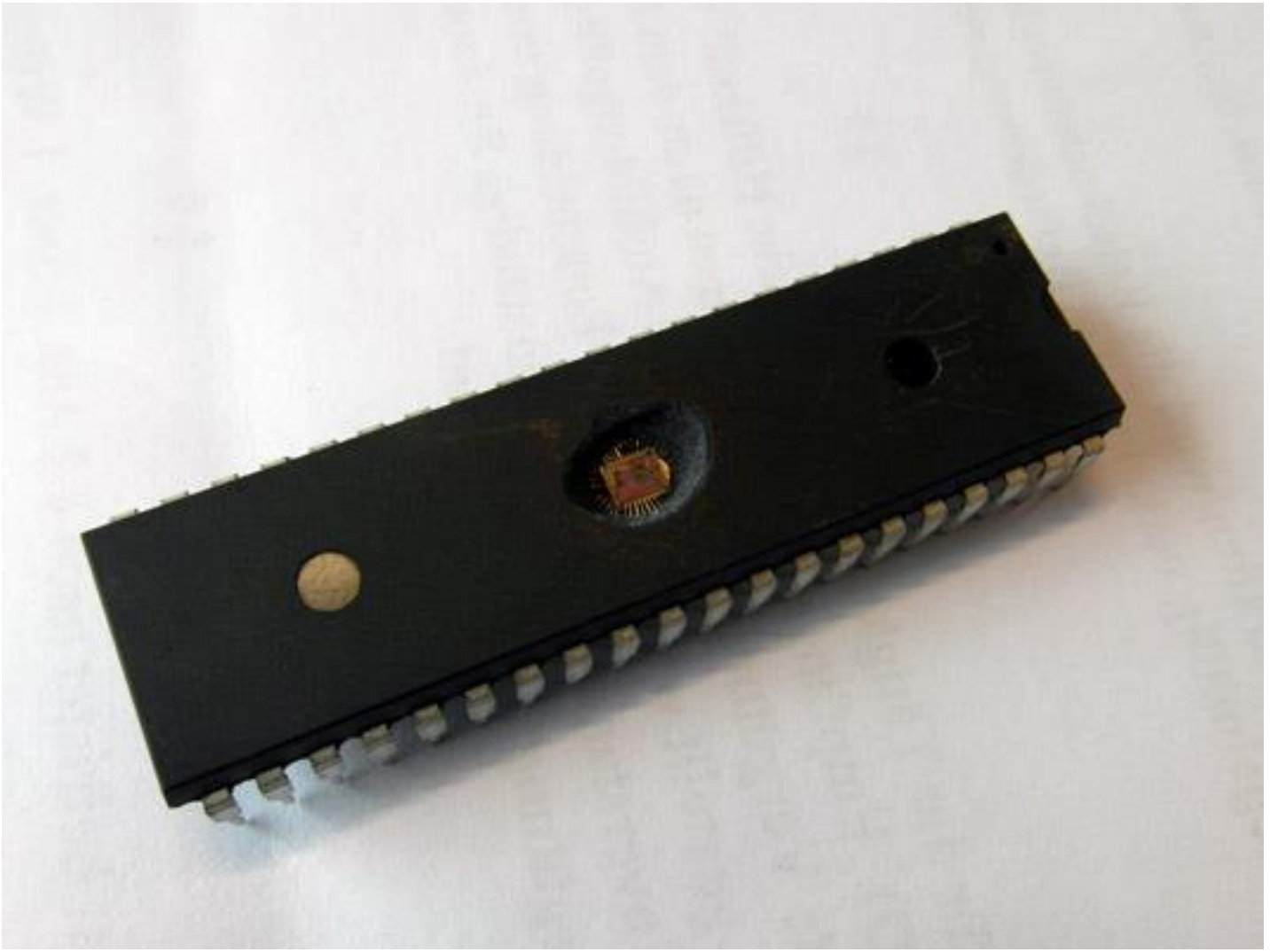


Decapping an IC (2)

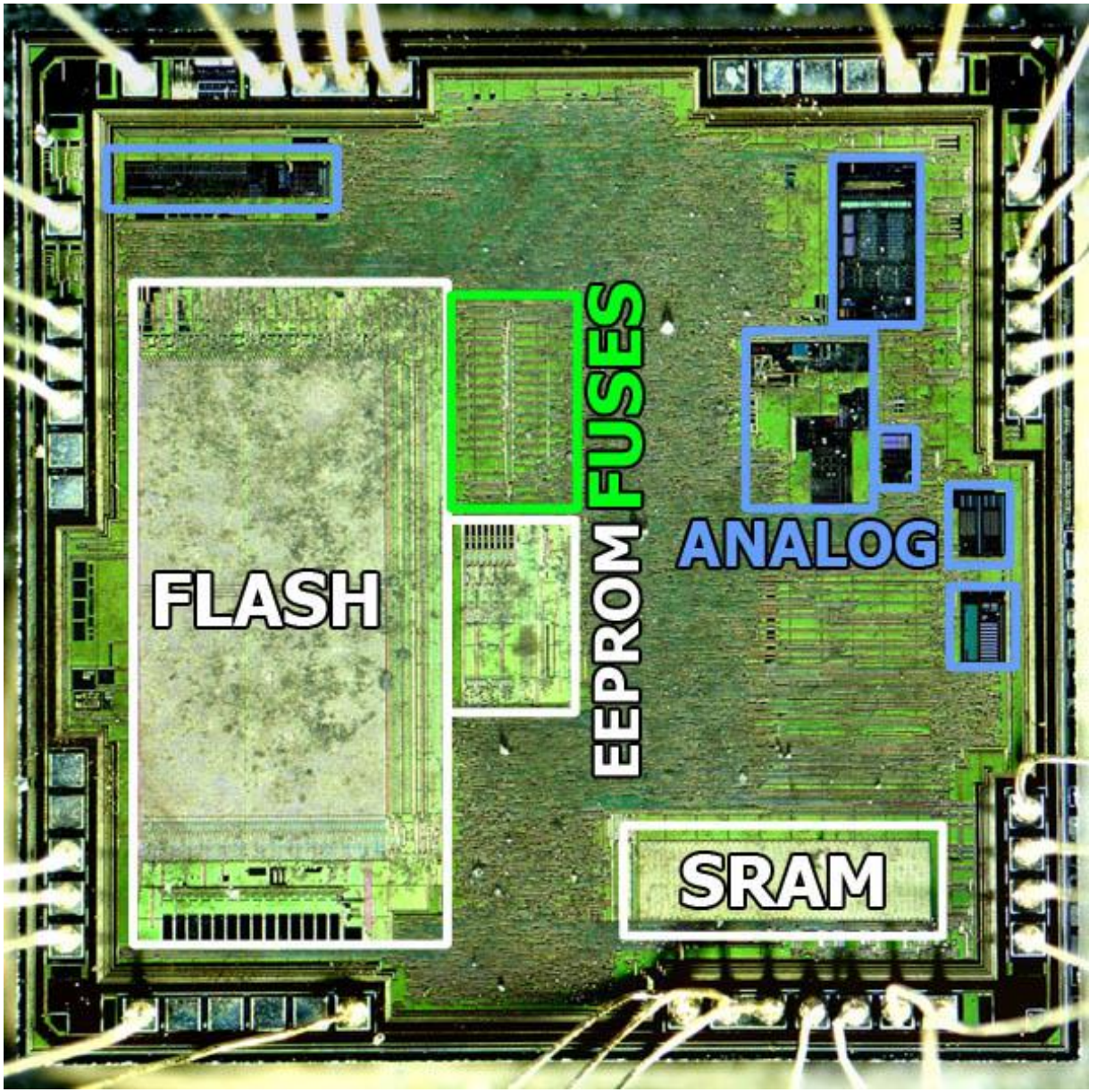


Decapping an IC (3)

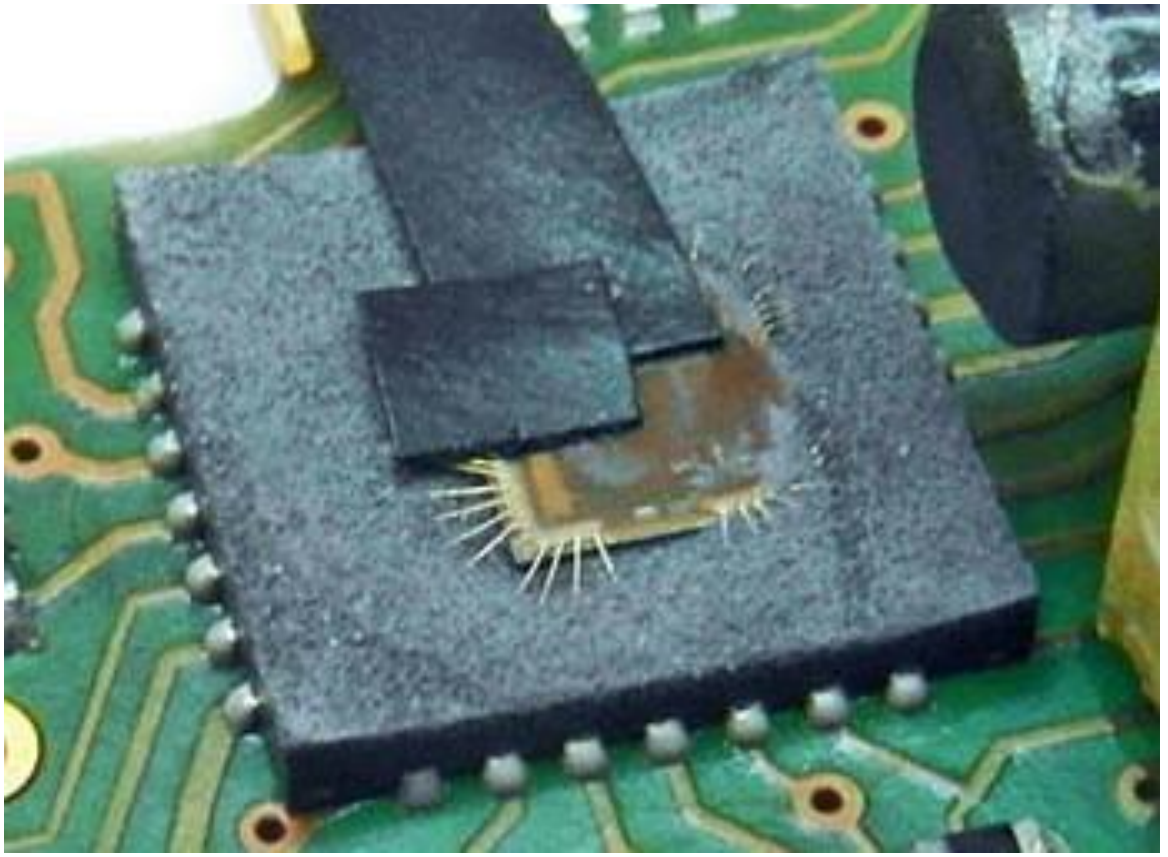
Decapping an IC (4)



Microscopic View of the Silicon Die



Exposure to UV-C: Disable Read-Out Protection (1)



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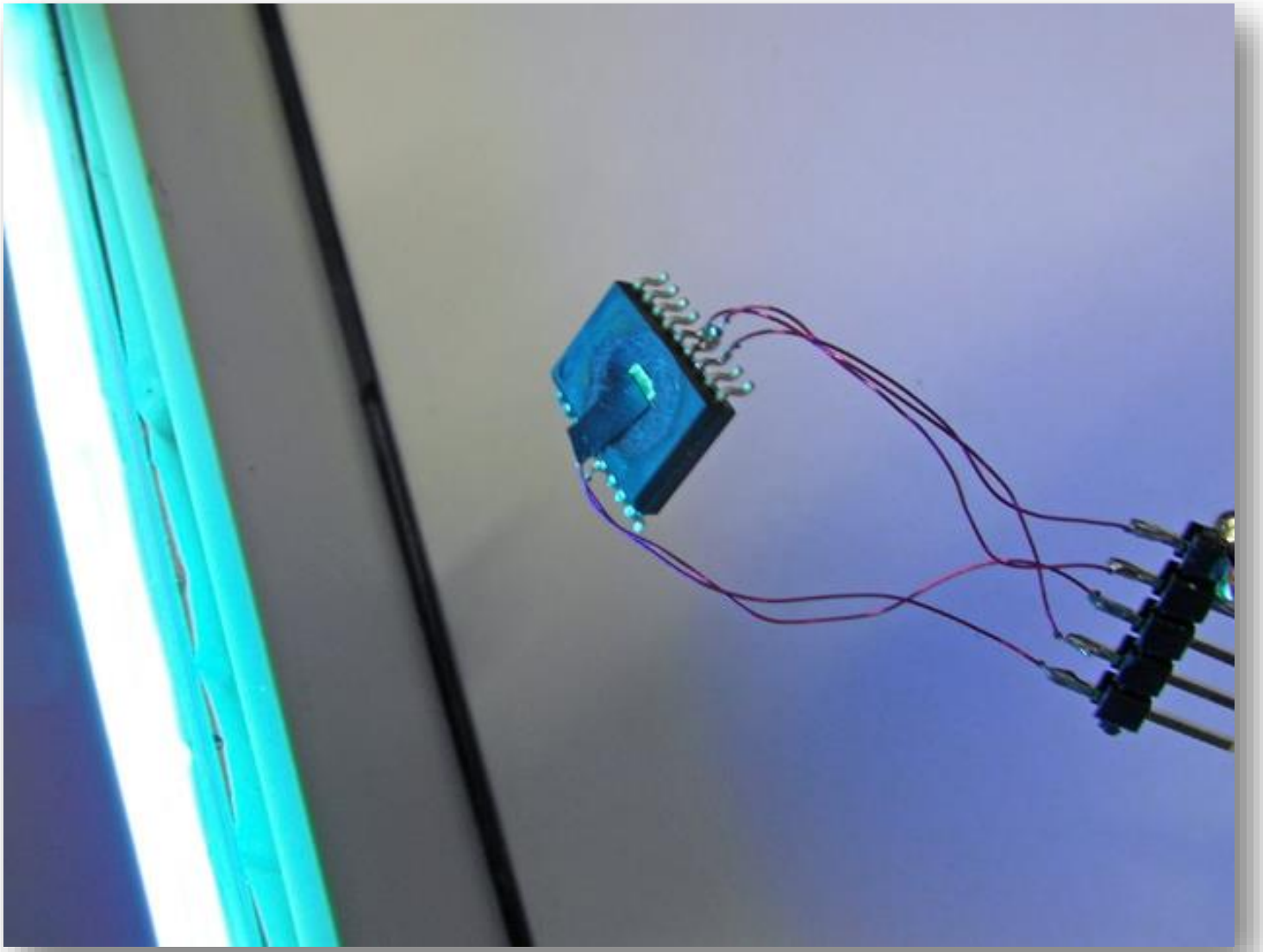
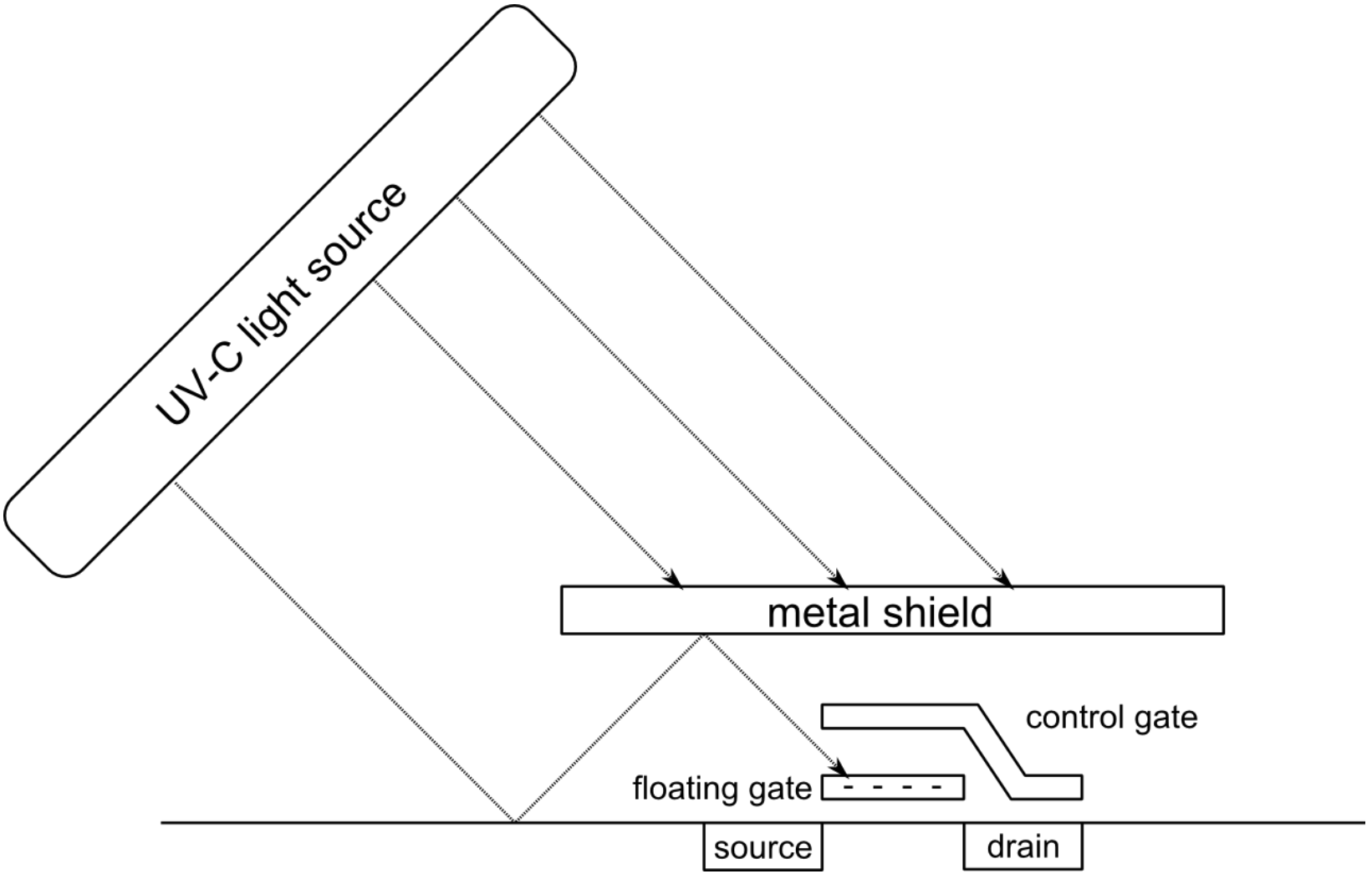
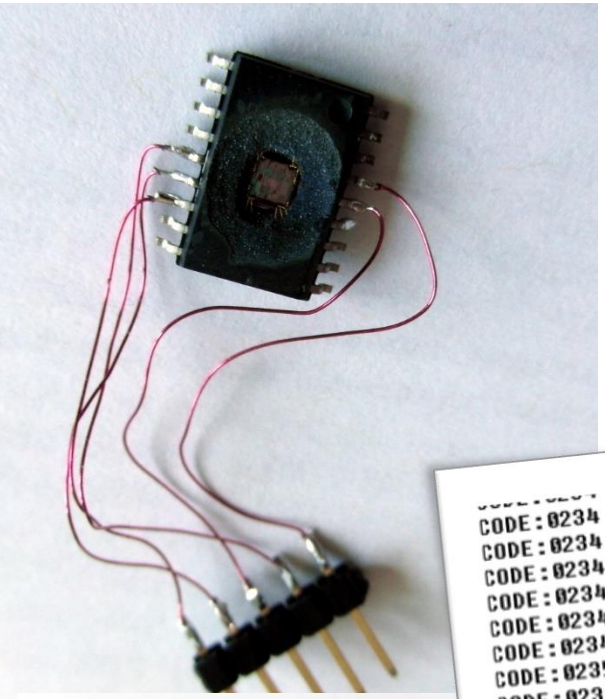


Exposure to UV-C: Why it works



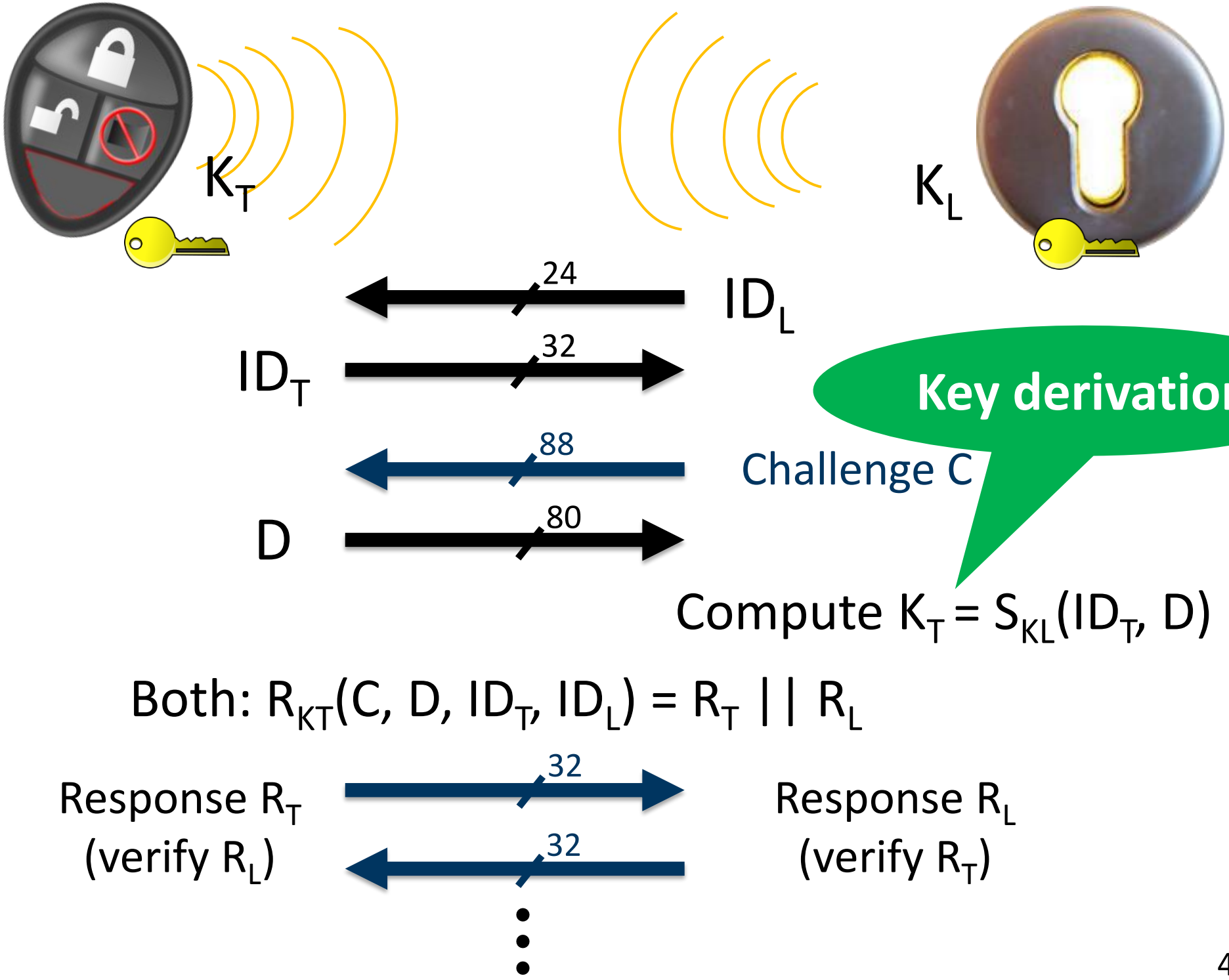
Reverse-Engineering continued



- Use standard programmer
- Reverse-Engineer (e.g., IDA)

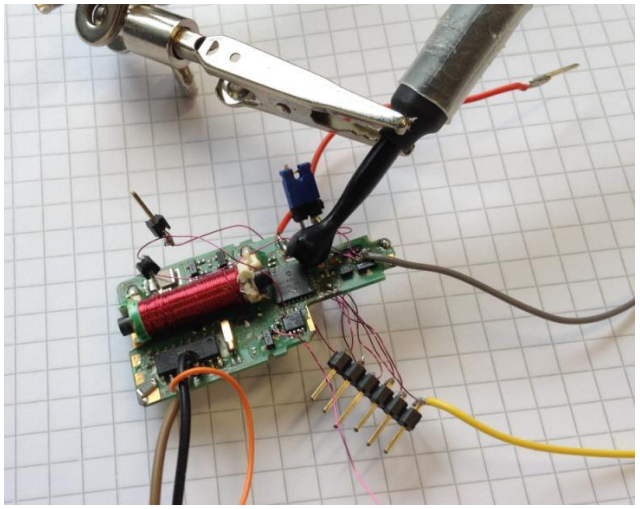
→ **all** internals known

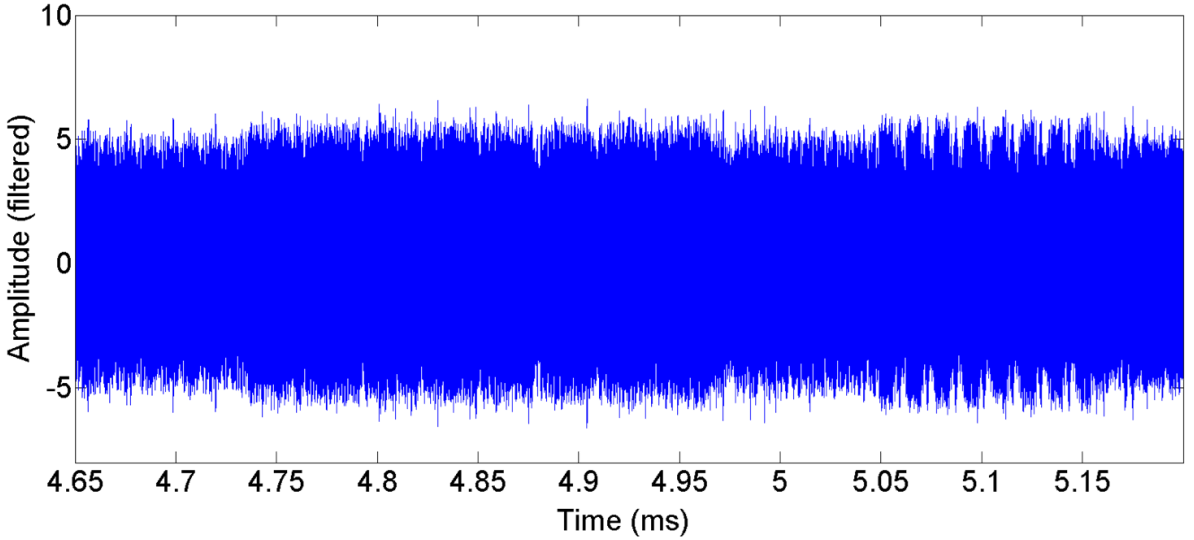
```
CODE:0234                                ; Init I2C
CODE:0234                                ; Set read address to Reg 71 / 70
CODE:0234                                ; Stores (inverted) read data at location pointed to by 73 / 7
CODE:0234                                ;
CODE:0234                                ; Sets (75) <- W
CODE:0234                                ; Sets bank to 0 (resets bit 5 and 6)
CODE:0235                        bsf     BANK0_STATUS, 6
CODE:0236 ; assume bank = 2
CODE:0236                        tstf    byte_DATA_119
CODE:0237                        skpnz
CODE:0238                                ; CODE XREF: maybe_related_to_rewriting_program_memory:endless
CODE:0238 endless_loop_1:
CODE:0238                        b       endless_loop_1
CODE:0239                        movlw   2
CODE:023A                        call    read_value_from_eeprom ; Adress in W
CODE:023A                                                       ; Result in W
CODE:023A                                                       ; Switches bank to 0
CODE:023B                        xorlw   0FF
CODE:023C                        bsf     BANK2_STATUS, 6
CODE:023D                        skpz
```

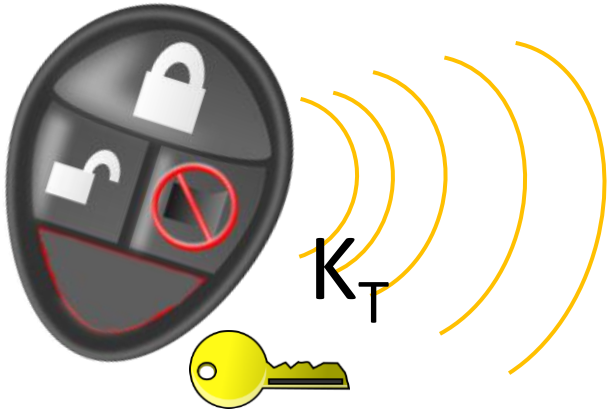

K_T
K_L
24
ID_L
ID_T
32
Key derivation
88
Challenge C
D
80
Compute $K_T = S_{KL}(ID_T, D)$
Both: $R_{KT}(C, D, ID_T, ID_L) = R_T \,||\, R_L$
Response R_T
(verify R_L)
32
32
Response R_L
(verify R_T)

Weaknesses and Attacks (1)

- Each lock stores installation-wide cryptographic key
- UV-C attack in ~ 30 min (decap PIC)
- Side-channel attack in ~ 15 min (access to PIC)
- **Attacking one lock gives access to all doors**







K_T

K_L

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ID_L

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Challenge C

Authentication

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Cryptographic Functions **R** and **S**

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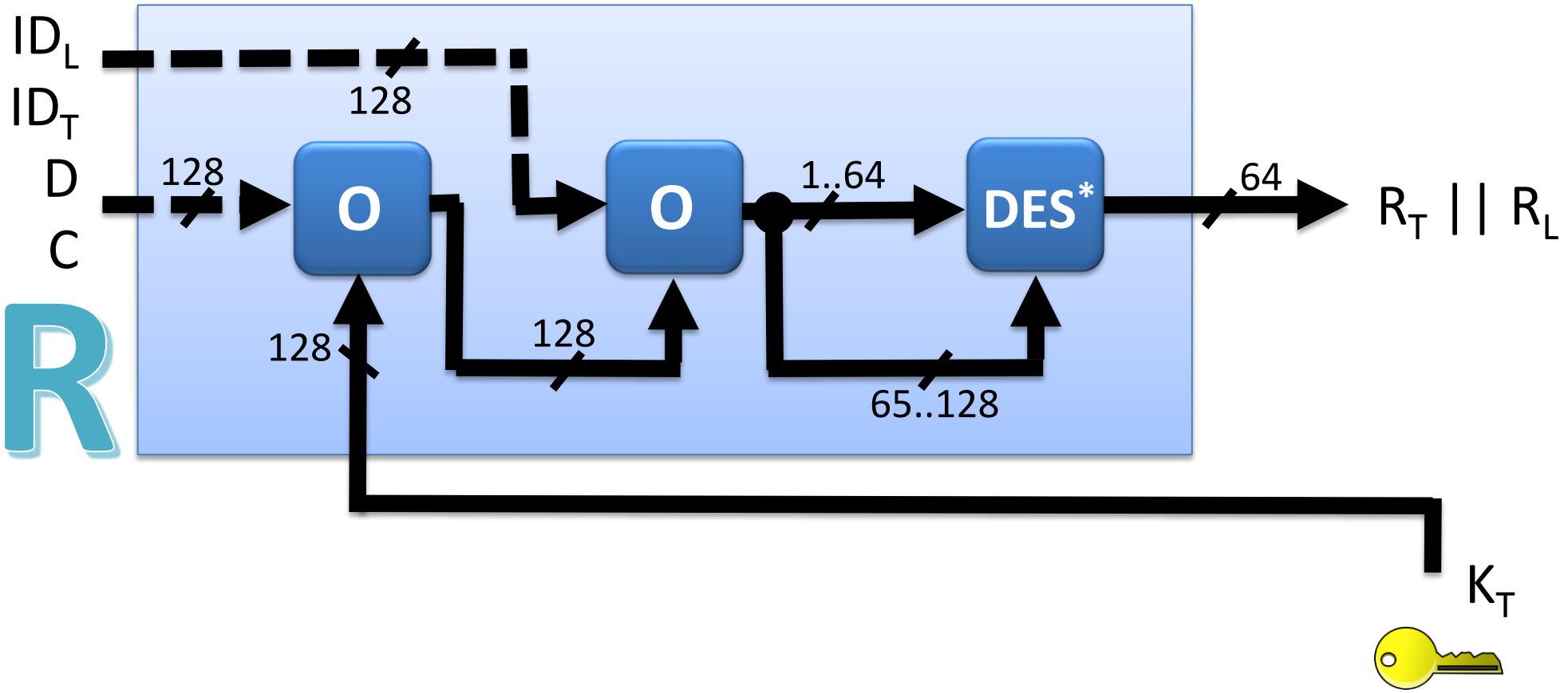
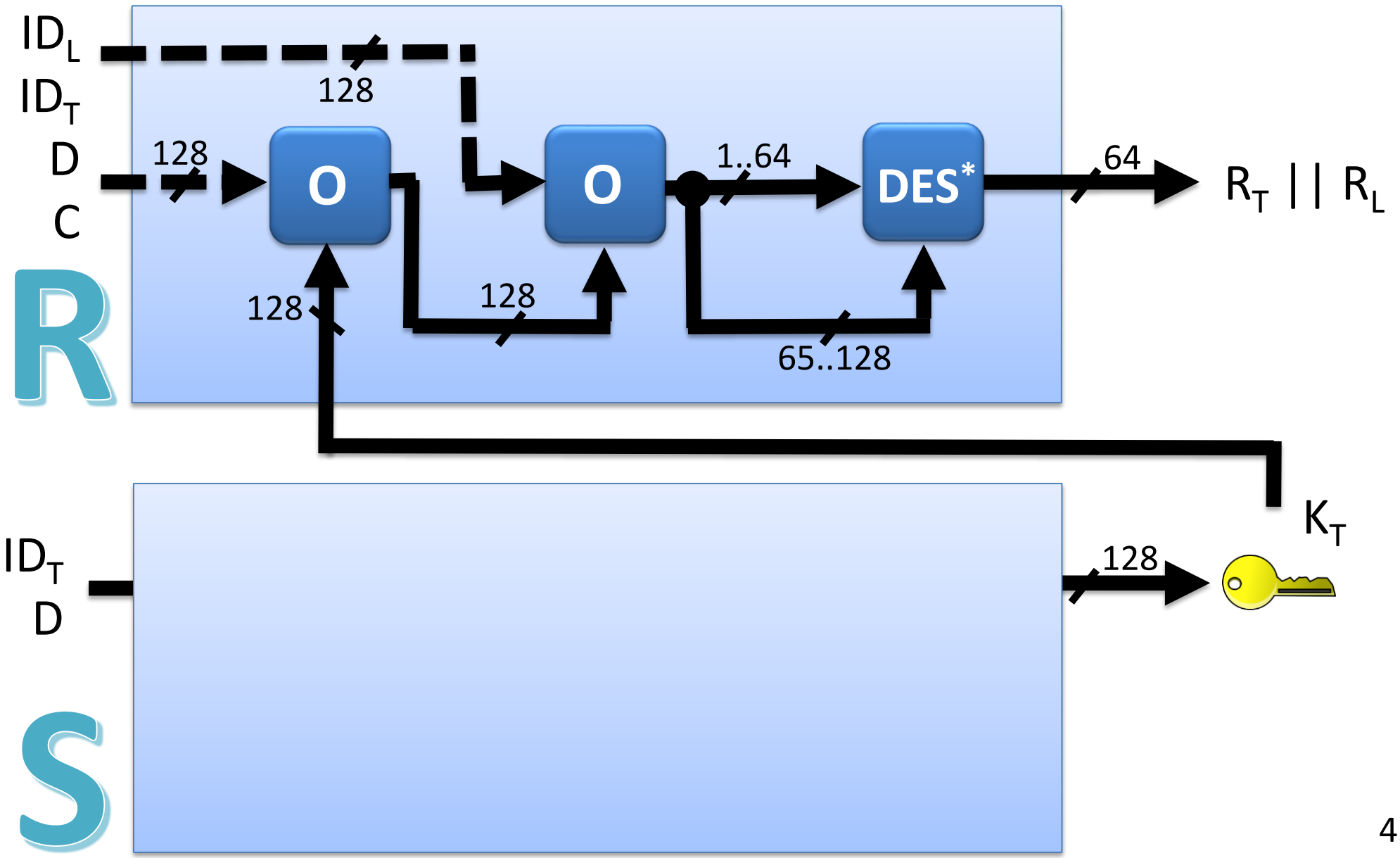
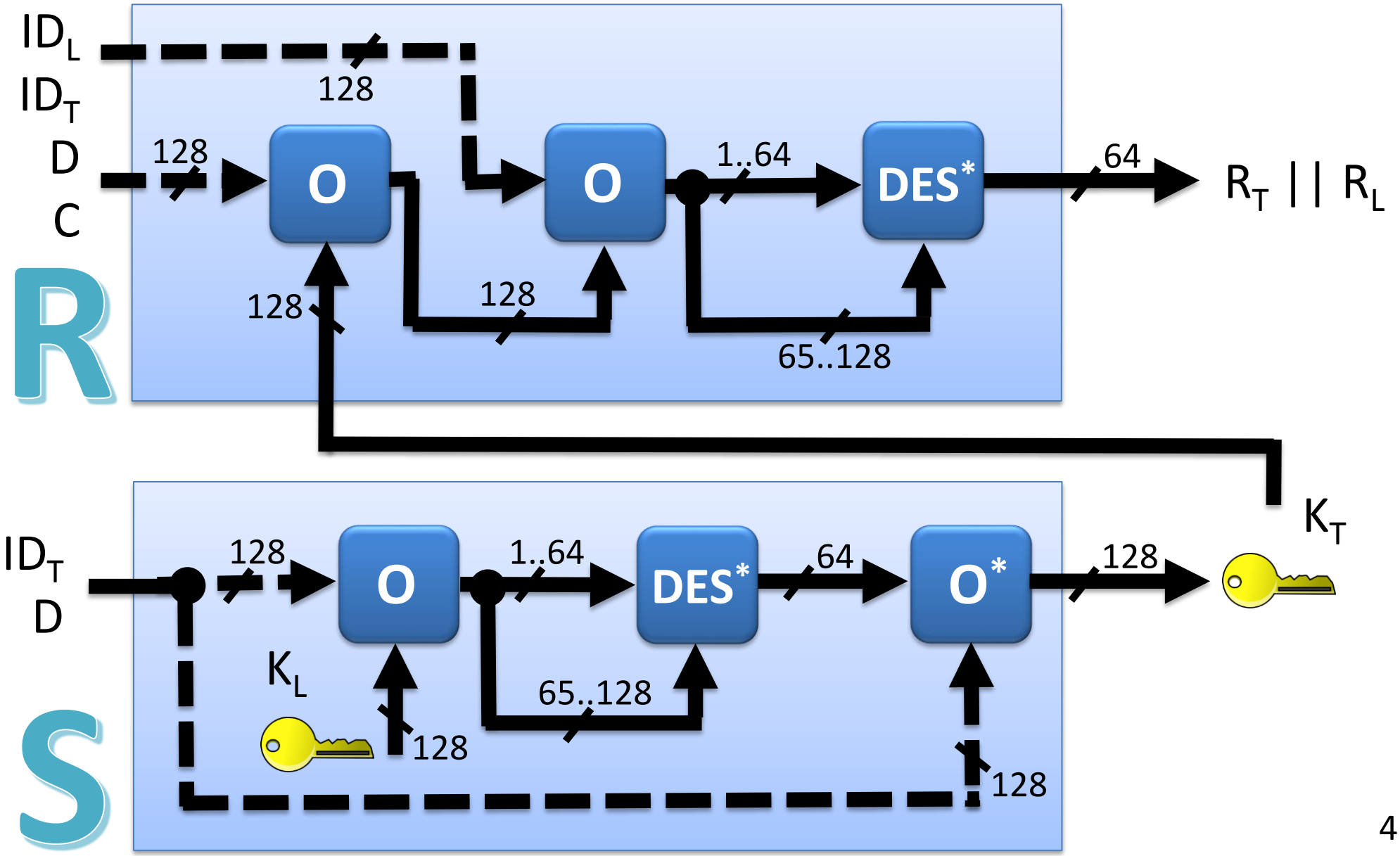


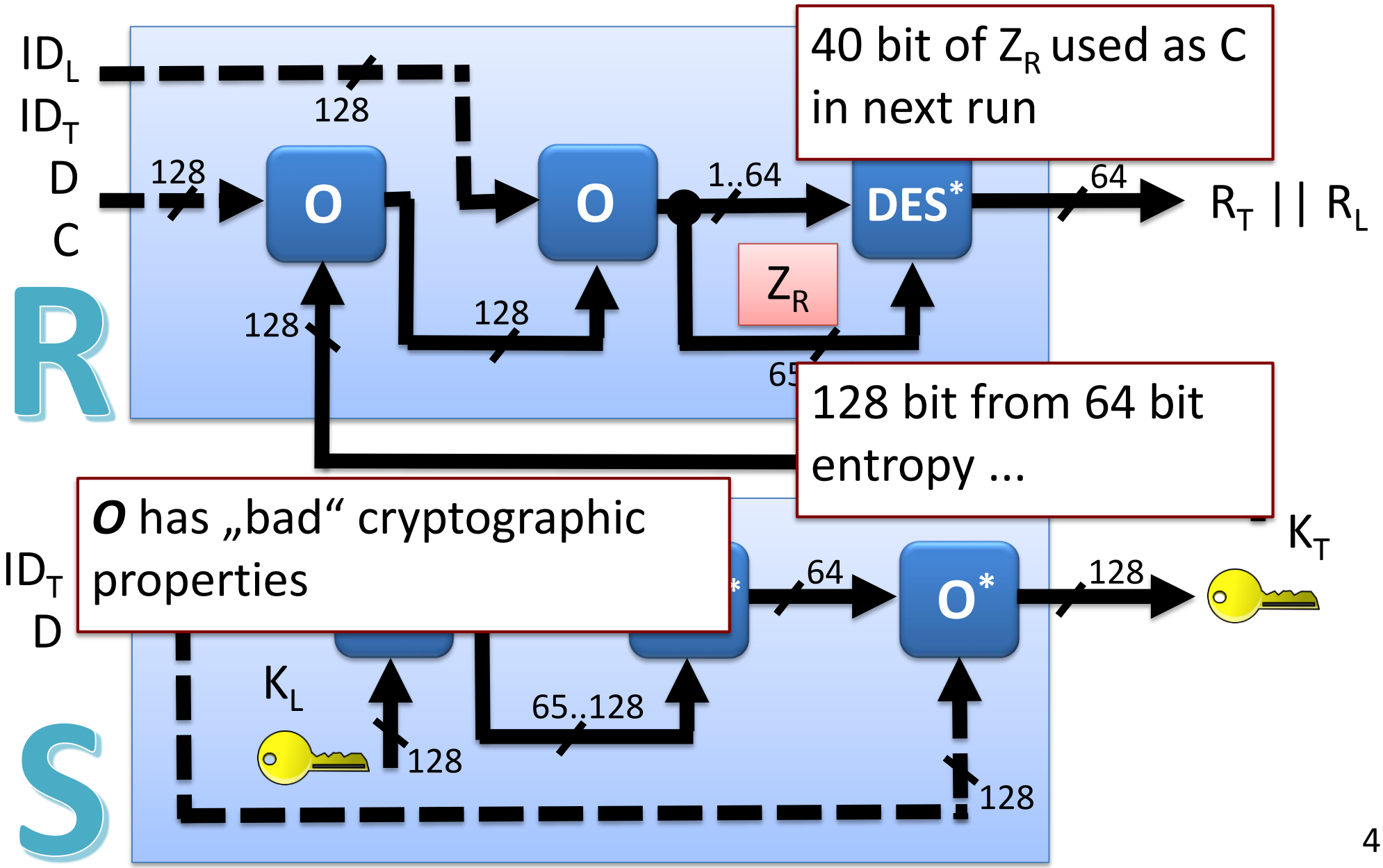
Cryptographic Functions **R** and **S**



Cryptographic Functions **R** and **S**



Cryptographic Functions **R** and **S**: Security Vulnerabilities



Consequence: Wireless **Lock-only Attack**

- Initate some, not successful protocol runs
- Compute K_T (for known ID_T)

Protocol Runs	Run-Time	Key Candidates
3	3,36 min	21,34
4	11,5 s	1
5	1,2 s	1
6	650 ms	1

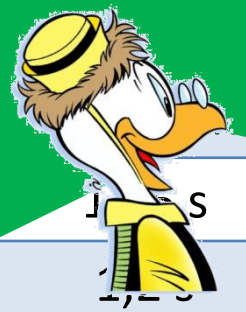
Consequence: Wireless **Lock-only Attack**

- Inita ns
- Com

Firmware upgrade of manufacturer: Generate **random C**

→ Fix for this attack

Protocol Runs		s
3		21,34
4	s	
5		1
6	650 ms	1



Report flaws

Improve

Level 3: Management Summary

- **Attacker can gain full access to any door**
- Reasons for security flaws
  - Insecure hardware
  - Proprietary cryptography
  - „Bad“ system design
- Can the system be „saved“?
  - **Cryptanalytical attacks:** Firmware update (cheap)
  - **HW attacks:** Require replacing all devices (expensive)

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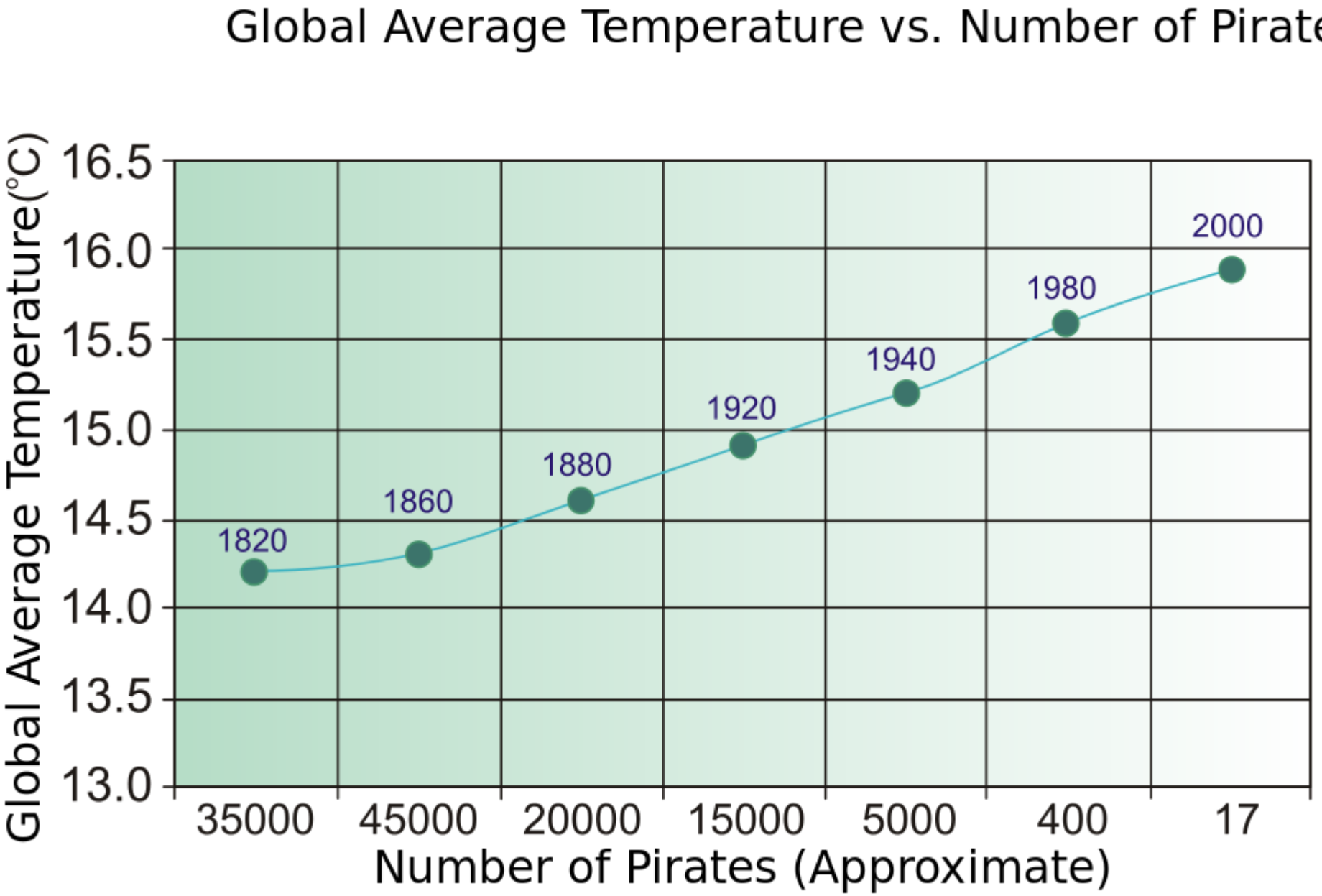
Responsible Disclosure

When pirates do good ...



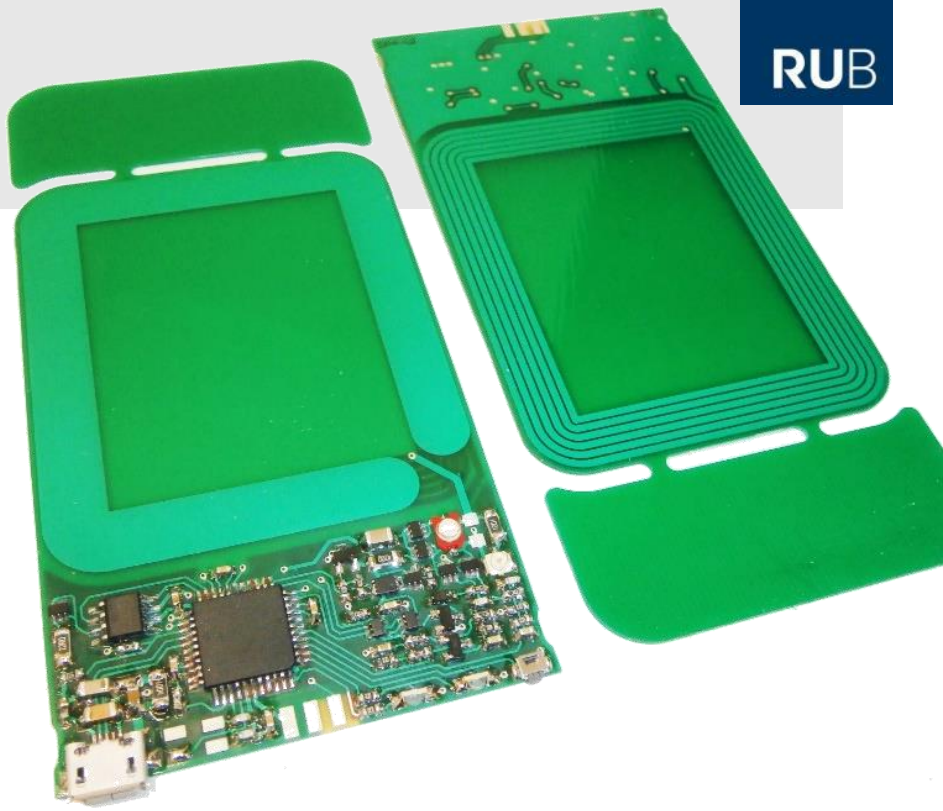
Responsible Disclosure

- **Locking system:**
  - Vendor informed ~ 1 year before
  - Discussion of found flaws
  - Deployed patch to fix mathematical attacks
- **Other examples:**
  - **Altera FPGAs:** Informed ~ 6 months before
  - **Yubikey**: Informed ~ 9 months before

Countermeasures

Countermeasures

- Implementation attacks: Practical threat, but:
- **First line of defense:** Classical countermeasures
  - Secure hardware (certified devices)
  - Algorithmic level
- **Second line of defense:** System level
  - **Detect**: Shadow accounts, logging
  - **Minimize impact** (where possible): Key diversification



Live Demo

„Everything that can go wrong, will go wrong“

Expect the unexpected.

Thanks!

Questions now?

or later:

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@sublevado